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Difficulties of malware analysis

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Safety from thinking



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Đơn vị tổ chức:



Đơn vị tài trợ:



Safety from thinking

Malware author perspective

- Manual Analysis
- Automated Malware Analysis

Malware author perspective

- Manual Analysis
 - How to confuse Malware Analyst ?
- Automated Malware Analysis
 - How to defeat Automated Malware Analysis ?



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Reverse Engineering approach

Safety from thinking

Analysis Prevention Techniques

- **Anti-Disassembly**
 - Protect malware from static analysis
- **Anti-Debugging**
 - Protect malware from dynamic analysis
- **Anti-VM**
 - Malware notice from VM environment
- **Anti-Sandbox**
 - Malware notice from sandbox environment



Anti-Disassembly

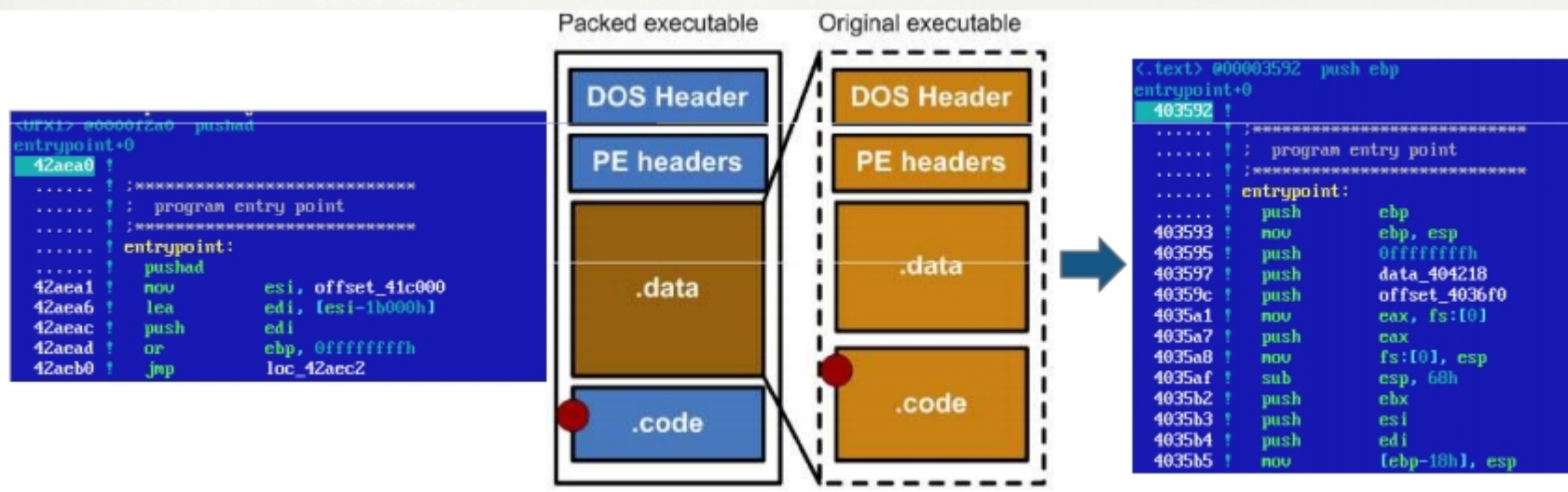
- Packing
- Confuse Disassembly Algorithms
- Fake Jumps
- Impossible disassembly
- ...



Packing

- Packers are used to shrink the size of executable file
- Makes static analysis harder

Packing





Packing

- Complex packers :
 - Multi-layer encryption
 - Advanced anti-debugging techniques
 - Code abstraction (metamorphic, virtual machines etc.)
 - Examples: Armadillo, Sdprotect, ExeCrypt, VMProtect



Packing

- Packer Analysis/Detection :
 - PEiD : it can detect more than 400 different signatures in PE files
 - RDG Packer Detector



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Packing

- More :
 - The Art of Unpacking
 - Anti-Unpacking Tricks - Peter Ferrie



Anti-Disassembly

- Type of Disassembly :
 - Linear sweep :
 - Disassemble one instruction at a time
 - Do not look at type of instruction
 - Recursive traversal :
 - Look at instruction and disassemble based on program flow
 - Used by IDA Pro and other commercial products



Anti-Disassembly

- Confuse Linear Disassembly Algorithm
 - Insert Garbage byte :
 - Since data is mixed with code, data can disassemble to valid instructions

```
jmp .destination
db 0x6a ; garbage byte
.destination:
; rest of the code
pop eax
```
 - Result of disassembler :

```
eb 01  jmp 0x401003
6a 58  push 0x58
```

Anti-Disassembly

- Confuse Linear Disassembly Algorithm
 - Insert Garbage byte :

PUSH ebp

MOV ebp, esp

XOR eax, eax



Anti-Disassembly

- Confuse Linear Disassembly Algorithm
 - Insert Garbage byte :

PUSH ebp

MOV ebp, esp

XOR eax, eax



CALL 31E58955





Anti-Disassembly

- Confuse Recursive traversal Disassembly Algorithm

Conditional that is, say, always true

- Jump Instructions with the Same Target

- The most common anti-disassembly technique seen in the wild is two back-to-back conditional jump instructions that both point to the same target

```
jz short near ptr loc_4011C4+1
```

```
75 01 jnz short near ptr loc_4011C4+1
```

- A Jump Instruction with a Constant Condition

- XOR : XOR instruction immediately followed by JNZ or JZ instruction

```
xor eax, eax
```

```
jz short near ptr loc_4011C4+1
```

- STC : STC instruction immediately followed by JNC or JAE instruction

- CLC : CLC instruction immediately followed by JC or JB instruction



Anti-Disassembly

- Splicing Instructions - Called "impossible disassembly"
 - A problem of representation

jmp -1

;these are hidden

inc eax

dec eax

Anti-Disassembly

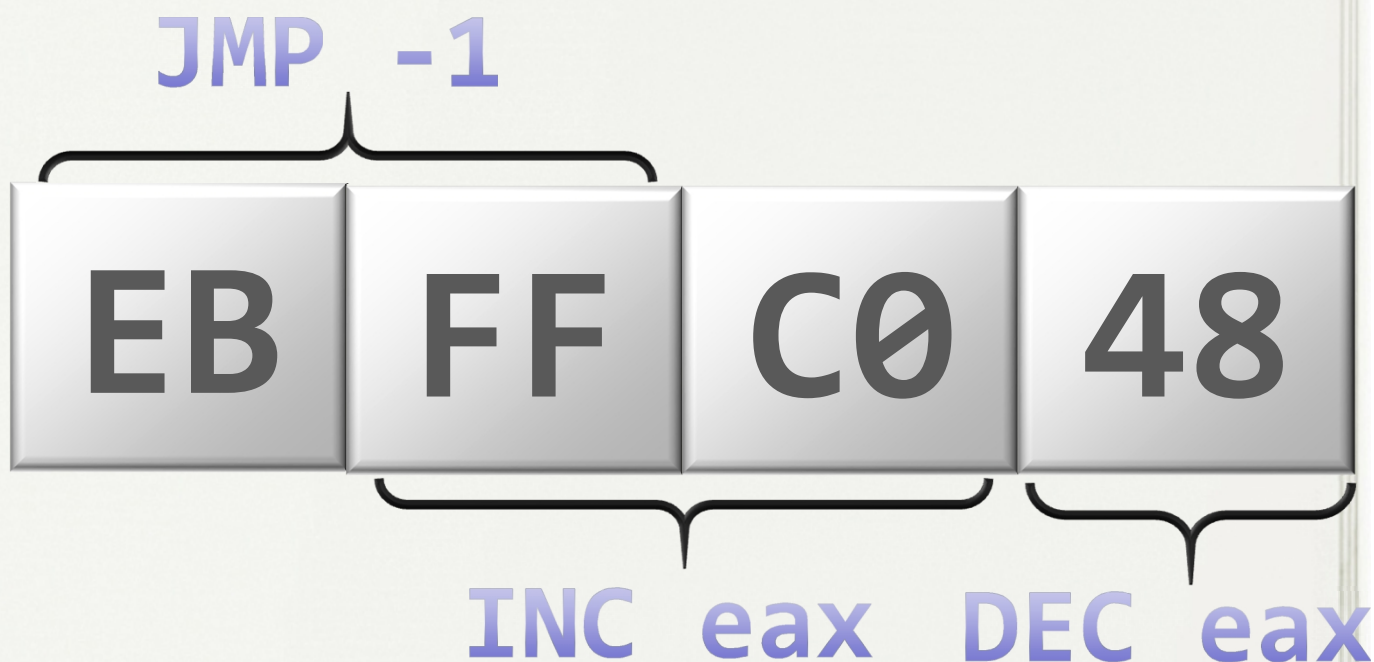
- Splicing Instructions - Called “impossible disassembly”
 - A problem of representation

jmp -1

;these are hidden

inc eax

dec eax





Anti-Disassembly

- Function pointer problems
 - It is easy to hide function calls made through pointers

- RET

004011C0	var_4	= byte ptr -4
004011C0	call	\$+5
004011C5	add	[esp+4+var_4], 5
004011C9	ret	
004011C9	sp-analysis	failed
004011CA	Confused IDA Pro	

- Misusing Structured Exception Handlers



Anti-Disassembly

- More :
 - Practical Malware analysis – chapter 15
 - <http://leetmatrix.blogspot.com/2013/02/an-anti-disassembly-trick.html>

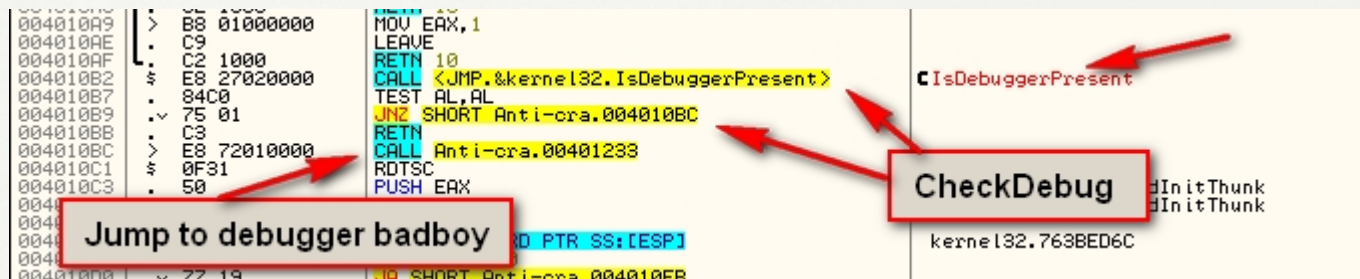


Anti-Anti-Disassembly

- IDA supports manually re-classifying code as well as code replacement to “fix” problem areas
- Deobfuscator : Deobfuscation plugin for IDA - <http://code.google.com/p/optimice/>
- Good malware analysts can recognize impossible assembly and run through the code to figure out what is going on

Anti-Debugging

- `IsDebuggerPresent()` Windows API



- if (IsDebuggerPresent()) {
- MessageBox(NULL, L"Debugger Detected Via IsDebuggerPresent",
- L"Debugger Detected", MB_OK);
- }



Anti-Debugging

- CheckRemoteDebuggerPresent() Windows API
 - CheckRemoteDebuggerPresent(GetCurrentProcess(), &pbIsPresent);
 - if (pbIsPresent) {
 - MessageBox(NULL, L"Debugger Detected Via
 - CheckRemoteDebuggerPresent", L"Debugger Detected", MB_OK);
 - }



Anti-Debugging

- IsDebuggerPresent : check the PEB.BeingDebugged flag
- ```
status = (_NtQueryInformationProcess) (hnd, ProcessBasicInformation, &pPIB, sizeof(PROCESS_BASIC_INFORMATION), &bytesWritten);
```
- ```
if (status == 0 ) {
```
- ```
 if (pPIB.PebBaseAddress->BeingDebugged == 1) {
```
- ```
        MessageBox(NULL, L"Debugger Detected Using PEB!IsDebugged", L"Debugger Detected", MB_OK);
```
- ```
 } else {
```
- ```
        MessageBox(NULL, L"No Debugger Detected", L"No Debugger Detected", MB_OK)
```



Anti-Debugging

- PEB ProcessHeap Flag Debugger Detection

```
• int main(int argc, char* argv[])
• {
•     unsigned int var;
•     __asm
•     {
•         MOV EAX, FS:[0x30];
•         MOV EAX, [EAX + 0x18];
•         MOV EAX, [EAX + 0x0c];
•         MOV var,EAX
•     }
•
•     if(var != 2)
•     {
•         printf("Debugger Detected");
•     }
•     return 0;
• }
```

Anti-Debugging

```
0:000> dt _PEB
ntdll!_PEB
+0x000 InheritedAddressSpace : UChar
+0x001 ReadImageFileExecOptions : UChar
+0x002 BeingDebugged : UChar
+0x003 SpareBool : UChar
+0x004 Mutant : Ptr32 Void
+0x008 ImageBaseAddress : Ptr32 Void
+0x00c Ldr : Ptr32 _PEB_LDR_DATA
+0x010 ProcessParameters : Ptr32 _RTL_USER_PROCESS_PARAMETERS
+0x014 SubSystemData : Ptr32 Void
+0x018 ProcessHeap : Ptr32 Void
+0x01c FastPebLock : Ptr32 _RTL_CRITICAL_SECTION
+0x020 FastPebLockRoutine : Ptr32 Void
+0x024 FastPebUnlockRoutine : Ptr32 Void
+0x028 EnvironmentUpdateCount : Uint4B
+0x02c KernelCallbackTable : Ptr32 Void
+0x030 SystemReserved : [1] Uint4B
+0x034 AtlThunkSListPtr32 : Uint4B
+0x038 FreeList : Ptr32 _PEB_FREE_BLOCK
+0x03c TlsExpansionCounter : Uint4B
+0x040 TlsBitmap : Ptr32 Void
+0x044 TlsBitmapBits : [2] Uint4B
+0x04c ReadOnlySharedMemoryBase : Ptr32 Void
+0x050 ReadOnlySharedMemoryHeap : Ptr32 Void
+0x054 ReadOnlyStaticServerData : Ptr32 Ptr32 Void
+0x058 AnsiCodePageData : Ptr32 Void
+0x05c OemCodePageData : Ptr32 Void
+0x060 UnicodeCaseTableData : Ptr32 Void
+0x064 NumberOfProcessors : Uint4B
+0x068 NtGlobalFlag : Uint4B
+0x070 CriticalSectionTimeout : _LARGE_INTEGER
+0x078 HeapSegmentReserve : Uint4B
+0x07c HeapSegmentCommit : Uint4B
+0x080 HeapDeCommitTotalFreeThreshold : Uint4B
+0x084 HeapDeCommitFreeBlockThreshold : Uint4B
+0x088 NumberOfHeaps : Uint4B
+0x08c MaximumNumberOfHeaps : Uint4B
+0x090 ProcessHeaps : Ptr32 Ptr32 Void
```



Anti-Debugging

- PEB!NtGlobalFlag
- ```
status = (_NtQueryInformationProcess) (hnd, ProcessBasicInformation, &pPIB, sizeof(PROCESS_BASIC_INFORMATION), &bytesWritten);
```
- ```
value = (pPIB.PebBaseAddress);
```
- ```
value = value+0x68;
```
- ```
if (*value == 0x70) {
```
- ```
 MessageBox(NULL, L"Debugger Detected Using PEB!NTGlobalFlag", L"Debugger Detected Using PEB!NTGlobalFlag", L"Debugger Detected", MB_OK);
```
- ```
} else {
```
- ```
 MessageBox(NULL, L"No Debugger Detected", L"No Debugger Detected", MB_OK);
```
- ```
}
```



Anti-Debugging

- RDTSC is used to retrieve the time stamp counter (number of clocks since boot-up) so this is a time-related trick. When you debug, the distance between those values that are returned in EAX will be higher than those when the program runs without being debugged. So, if there is really a difference you're debugging

Anti-Debugging

00401078	75 2F	JNZ SHORT Anti-cra.004010A9		
0040107A	66 3D B90B	CMP AX,0BB9		
0040107E	75 1E	JNZ SHORT Anti-cra.0040109F		
00401080	8B 45 08	MOV EAX,ARG_1		
00401083	A3 78304000	MOV DWORD PTR DS:[403078],EAX		
00401088	33 C0	XOR EAX,EAX		
0040108A	A3 86304000	MOV DWORD PTR DS:[403086],EAX		
0040108F	E8 2D000000	CALL Anti-cra.004010C1		
00401094	E8 57000000	CALL Anti-cra.004010F0		
00401099	E8 72000000	CALL Anti-cra.00401110		
0040109E	EB 09	JMP SHORT Anti-cra.004010A9		
004010A0	B8 00000000	MOV EAX,0		
004010A5	C9	LEAVE		
004010A6	C2 1000	RETN 10		
004010A9	B8 01000000	MOV EAX,1		
004010AE	C9	LEAVE		
004010AF	C2 1000	RETN 10		
004010B2	E8 27020000	CALL <JMP.&kernel32.IsDebuggerPresent>	IsDebuggerPresent	
004010B7	84 C0	TEST AL,AL		
004010B9	90	NOP		
004010BA	90	NOP		
004010BB	C3	RETN		
004010BC	E8	DB E8		
004010BD	72	DB 72		
004010BE	01	DB 01		
004010BF	00	DB 00		
004010C0	00	DB 00		
004010C1	50	PUSH EAX		
004010C3	33 C0	XOR EAX,EAX		
004010C4	0F 31	RDTSC		
004010C6	2B 0424	SUB EAX,DWORD PTR SS:[ESP]		
004010C8	3D 00200000	CMP EAX,2000		
004010CB	77 19	JA SHORT Anti-cra.004010EB		
004010D2	6A 0A	PUSH 0A		

Registers (FPU)

EAX	4E763511
ECX	00000000
EDX	00000FD7
EBX	00000000
ESP	0012FAA0
EBP	0012FAA4
ESI	00000111
EDI	0012FB20
EIP	004010C3 Anti-cra.0
C 0	ES 0023 32bit 0(FF
P 1	CS 001B 32bit 0(FF
A 0	SS 0023 32bit 0(FF
Z 1	DS 0023 32bit 0(FF
S 0	FS 003B 32bit 7FFD
T 0	GS 0000 NULL
D 0	
O 0	LastErr ERROR_SUCC
EFL	00000246 (NO,NB,E,B
ST0	empty 0.0
ST1	empty 0.0
ST2	empty 0.0
ST3	empty 0.0
ST4	empty 0.0
ST5	empty 0.0
ST6	empty 0.0
ST7	empty 0.0
FST	4020 Cond 1 0 0 0
FCW	027F Prec NEAR,53

Our breakpoint

Call to CheckTiming

First RDTSC

EAX = current system time

004010C3	50	PUSH EAX		
004010C4	33 C0	XOR EAX,EAX		
004010C6	0F 31	RDTSC		
004010C8	2B 0424	SUB EAX,DWORD PTR SS:[ESP]		
004010CB	3D 00200000	CMP EAX,2000		
004010D0	77 19	JA SHORT Anti-cra.004010EB		
004010D2	6A 0A	PUSH 0A		
004010D4	68 7C304000	PUSH Anti-cra.0040307C		
004010D9	68 B80B0000	PUSH 0BB8		
004010DE	FF 35 78304000	PUSH DWORD PTR DS:[403078]		

MUCH larger than 2000h

Registers (FPU)

EAX	E0557554
ECX	00000000
EDX	0000106C
EBX	00000000
ESP	0012FA9C
EBP	0012FAA4
ESI	00000111
EDI	0012FB20
EIP	004010D0

Count = A (10.)
Buffer = Anti-cra.0040307C
ControlID = BB8 (3000...
hWnd = 0010058E ('Anti-



Anti-Debugging

```
• int main(int argc, char* argv[])
• {
•     unsigned int time1 = 0;
•     unsigned int time2 = 0;
•     __asm
•     {
•         RDTSC
•         MOV time1,EAX
•         RDTSC
•         MOV time2, EAX
•
•     }
•     if ((time2 - time1) > 100)
•     {
•         printf("%s", "VM Detected");
•         return 0;
•     }
•     printf("%s", "VM not present");
•     return 0;
• }
```



Anti-Debugging

- Find evidence of debugger on system:
 - Registry entries
 - FindWindow API call :
 - `HANDLE ollyHandle = NULL;`
 - `ollyHandle = FindWindow(L"OLLYDBG", 0);`
 - `if (ollyHandle == NULL) {`
 - `MessageBox(NULL, L"OllyDbg Not Detected", L"Not Detected", MB_OK);`
 - `} else {`
 - `MessageBox(NULL, L"Ollydbg Detected Via OllyDbg FindWindow()", MessageBox(NULL, L"Ollydbg Detected Via OllyDbg FindWindow()",`
 - `L"OllyDbg Detected", MB_OK);`
 - `}`

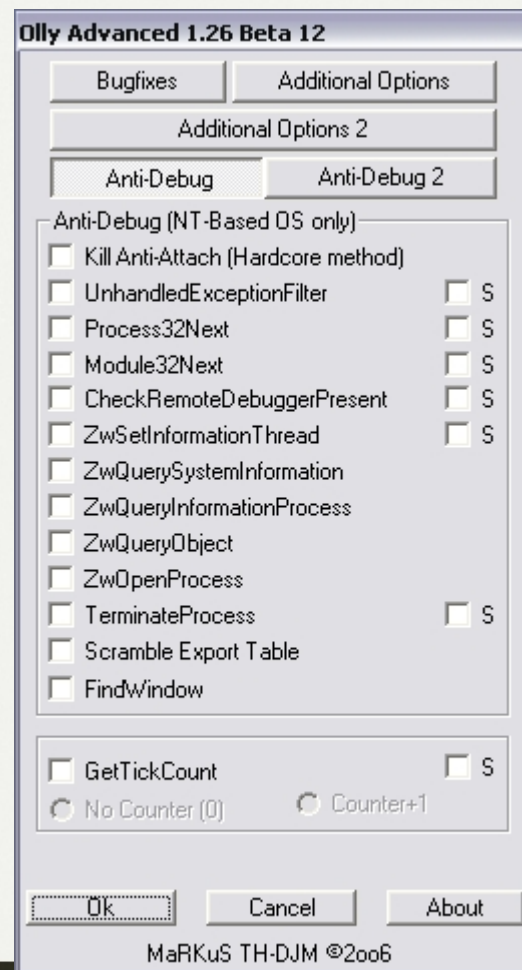


Anti-Debugging

- More :
 - Anti-Debugging - A Developers Viewpoint
 - Windows Anti-Debug Reference
 - The “Ultimate” Anti-Debugging Reference

Anti-anti-debugging

- Olly Advanced
- StrongOD
- aadp : <http://code.google.com/p/aadp/>





Anti-VM

- VM Fingerprints
 - Descriptor Table addresses (IDT, LDT, etc.)
 - Running Processes (eg. VMWare Tools)
 - Registry entries that include "VMWare"
 - loaded modules name
 - Default virtual machine hardware
 - Common VM MAC addresses
 - VMWare specific I/O port
 - Basically, any difference between a VM and a real computer

Anti-VM

```
00428890  
00428890 loc_428890:  
00428890 push    dword ptr [eax]  
00428892 call    hashString_0  
00428897 pop     ecx  
00428898 cmp     eax, 2FE483F3h ; vm SCSI.sys  
0042889D jz      short loc_4288F2  
  
0042889F cmp     eax, 2FD5F8F3h ; vm Hgfs.sys  
004288A4 jz      short loc_4288F2  
  
004288A6 cmp     eax, 0CFF129A8h ; vmx Sgxa.sys  
004288AB jz      short loc_4288F2  
  
004288AD cmp     eax, 2FDB60F3h ; vmxnet.sys  
004288B2 jz      short loc_4288F2  
  
004288B4 cmp     eax, 2FFC2FB4h ; vm mouse.sys  
004288B9 jz      short loc_4288F2  
  
004288BB cmp     eax, 2FF91C94h ; vm debug.sys  
004288C0 jz      short loc_4288F2
```

Anti-VM

```
004288C8  
004288C8 loc_4288C8:                ; vmwaretray  
004288C8 push    9F408DC4h  
004288CD call    enumCheckProcesses ; result in EAX...EAX=0 means passed  
004288D2 pop     ecx  
004288D3 test   eax, eax  
004288D5 jnz    short End  
  
004288D7 push    9F5784C4h        ; vmwareuser  
004288DC call    enumCheckProcesses ; result in EAX...EAX=0 means passed  
004288E1 pop     ecx  
004288E2 test   eax, eax  
004288E4 jnz    short End
```



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Anti-VM

```
• int main(int argc, char **argv)
• {
•
• char lszValue[100];
• HKEY hKey;
• int i=0;
• RegOpenKeyEx (HKEY_LOCAL_MACHINE, "SYSTEM\\CurrentControlSet\\Services\\Disk\\Enum", 0L,
KEY_READ , &hKey);
•
• RegQueryValue(hKey,"0",lszValue,sizeof(lszValue));
•
• printf("%s", lszValue);
• if (strstr(lszValue, "VMware"))
• {
•     printf("Vmware Detected");
• }
•
• RegCloseKey(hKey);
• return 0;
•
```



Anti-VM

- Red Pill is an anti-VM technique that executes the sidt instruction to grab the value of the IDTR register. The virtual machine monitor must relocate the guest's IDTR to avoid conflict with the host's IDTR. Since the virtual machine monitor is not notified when the virtual machine runs the sidt instruction, the IDTR for the virtual machine is returned. The Red Pill tests for this discrepancy to detect the usage of VMware.



Anti-VM

- The sgdt and sltd instruction technique for VMware detection is commonly known as No Pill. Unlike Red Pill, No Pill relies on the fact that the LDT structure is assigned to a processor, not an operating system. And because Windows does not normally use the LDT structure, but VMware provides virtual support for it, the table will differ predictably : The LDT location on the host machine will be zero, and on the virtual machine, it will be nonzero. A simple check for zero against the result of the sltd instruction does the trick.



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Anti-VM

- ScoopyNG - The VMware Detection Tool
 - ScoopyNG combines the detection tricks of Scoopy Doo and Jerry as well as some new techniques to determine if a current OS is running inside a VMware Virtual Machine (VM) or on a native system.
 - The first three checks look for the sidt, sgdt, and sltd (Red Pill and No Pill) instructions.
 - The fourth check looks for str.
 - The fifth and sixth use the backdoor I/O port 0xa and 0x14 options, respectively.
 - The seventh check relies on a bug in older VMware versions running in emulation mode.



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Anti-VM

```
C:\WINDOWS\system32\cmd.exe
#####
::      ScoopyNG - The VMware Detection Tool      ::
::      Windows version v1.0                      ::
::
[+] Test 1: IDT
IDT base: 0x82c89400
Result  : Native OS

[+] Test 2: LDT
LDT base: 0xdead0000
Result  : Native OS

[+] Test 3: GDT
GDT base: 0x82c89000
Result  : Native OS

[+] Test 4: STR
STR base: 0x28000000
Result  : Native OS

[+] Test 5: VMware "get version" command
Result  : Native OS

[+] Test 6: VMware "get memory size" command
Result  : Native OS

[+] Test 7: VMware emulation mode
Result  : Native OS or VMware without emulation mode
          (enabled acceleration)

::      tk, 2008                                  ::
::      [ www.trapkit.de ]                        ::
#####
```

```
#####
::      ScoopyNG - The VMware Detection Tool      ::
::      Windows version v1.0                      ::
::
[+] Test 1: IDT
IDT base: 0xffc18000
Result  : VMware detected

[+] Test 2: LDT
LDT base: 0xdead4060
Result  : VMware detected

[+] Test 3: GDT
GDT base: 0xffc07000
Result  : VMware detected

[+] Test 4: STR
STR base: 0x00400000
Result  : VMware detected

[+] Test 5: VMware "get version" command
Result  : VMware detected
Version : Workstation

[+] Test 6: VMware "get memory size" command
Result  : VMware detected

[+] Test 7: VMware emulation mode
Result  : Native OS or VMware without emulation mode
          (enabled acceleration)

::      tk, 2008                                  ::
::      [ www.trapkit.de ]                        ::
#####
```



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Anti-VM

- The same with VirtualBox
 - 9 method to detect VirtualBox by waleedassar : <http://pastebin.com/RU6A2UuB>



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Anti-VM

- More :
 - Thwarting Virtual Machine Detection
 - Detecting the Presence of Virtual Machines Using the Local Data Table



Anti-anti-vm

- Hardening your VM
 - Don't install the VMware tool
 - Change the configuration of your virtual machine by adding the following options to your .vmx file :
 - isolation.tools.getPtrLocation.disable = "TRUE"
 - isolation.tools.setPtrLocation.disable = "TRUE"
 - isolation.tools.setVersion.disable = "TRUE"
 - isolation.tools.getVersion.disable = "TRUE"
 - monitor_control.disable_directexec = "TRUE"
 - monitor_control.disable_chksimd = "TRUE"
 - monitor_control.disable_ntreloc = "TRUE"
 - monitor_control.disable_selfmod = "TRUE"
 - monitor_control.disable_reloc = "TRUE"
 - monitor_control.disable_btinout = "TRUE"
 - monitor_control.disable_btmemspace = "TRUE"
 - monitor_control.disable_btpriv = "TRUE"
 - monitor_control.disable_btseg = "TRUE"



Anti-anti-vm

- Patching the code : If you debug the malware and identify some of the specific instructions (e.g. *sidt*, *sgdt*, *sldt*) you can replace the code with NOPs to prevent it.
- More :
 - http://handlers.sans.org/tliston/ThwartingVMDetection_Liston_Skoudis.pdf
 - http://radlab.cs.berkeley.edu/w/upload/3/3d/Detecting_VM_Aware_Malware.pdf
 - <http://vrt-blog.snort.org/2009/10/how-does-malware-know-difference.html>



Anti-Sandbox

- Logic bombs are particular checks in the program which require certain events to be true in order to execute the malicious payload
 - Checking if something changes on the user desktop
 - checking if the mouse pointer is not moving for a particular time
- Sandbox Overloading
 - Malware flood the sandbox by generating too much worthless behavior data (e.g the sleep call) before executing the real payload. Logging the generated behavior data introduces additional delays and therefore the execution does not reach the real payload
 - Solution : only analysis network traffic, does not capture any system level behavior



Anti-Sandbox

- `#include <windows.h>`
`#include <stdio.h>`
- `void overloadSandbox()`
`{`
 `char Path[20];`

 `for(int i = 0; i < 100000; ++i)`
 `{`
 `sprintf_s(Path, 20, "C:\\test");`
 `DeleteFile(Path);`
 `}`
`}`


```
int main()
{
    overloadSandbox();
    //Real payload
    CreateFile("C:\\payload", GENERIC_WRITE, FILE_SHARE_WRITE, NULL,
    CREATE_ALWAYS, FILE_ATTRIBUTE_NORMAL, NULL);

    return 0;
}
```



Anti-Sandbox

- Pafish is a demo tool that performs some anti(debugger/VM/sandbox) tricks :
 - <https://github.com/a0rtega/pafish>



Anti-Anti-Sandbox

Logic bombs :

- By understanding the behavior of the logic bomb code in the analysis report human analysts, we can improve your sandbox

Sandbox Overload :

- In some case, we should only analyses only network traffic and does not capture any system level behavior
- In addition we could also write a signature to detect and blacklist the massive worthless behavior data (e.g the sleep call)



Some tools can help you

- Crowd Detox : plugin for Hex-Rays automatically removes junk code and variables from Hex-Rays function decompilations
- CrowdRE : aims to make it easier for developers to reverse engineer complex applications by working collaboratively with other users
- <http://www.crowdstrike.com/community-tools/index.html>



Conclusion

- Automated Malware Analysis is good but it can be defeat by new anti-* techniques => we still need manual analysis for advance malwares and update back to AMAs



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Thank you !